

UNIT 8

COMBINATIONAL CIRCUIT DESIGN AND SIMULATION



Fall 2021

Combinational Circuit Design & Simulation

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 - Design of circuits with limited gate fan-in
 - Gate delay and timing diagrams
 - Hazards in combinational logic
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 - Unit 8

Review of Combinational Circuit Design

- **Generic combinational circuit design steps:**
 1. Specify the switching function(s)
 - A truth table
 - Algebraic equations
 2. Derive simplified algebraic expression(s)
 - Karnaugh maps
 - Algebraic simplification (large # of variables, small # of terms)
 3. Manipulate into the proper form
 - Apply algebraic simplification with considering the gate types used in circuit implementation
- **Minimum SOP (POS) \Rightarrow minimum two-level circuits**
 - Start from min SOP \equiv AND-OR \Rightarrow NAND-NAND/OR-NAND/NOR-OR
 - Start from min POS \equiv OR-AND \Rightarrow NOR-NOR/AND-NOR/NAND-AND
- **The # of level = the max # of cascaded gates between I/Os**
 - Change level by factoring or multiplying out
 - \Rightarrow Change # of gates or gate inputs
- **Design a multi-level multi-output ckt from a ckt of AND & OR gates**

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Limited Gate Fan-in

For a gate, fan-in = # of inputs
For a circuit, fan-in = max # of inputs on each gate

Example 1

- Realize $f(a, b, c, d) = \Sigma m(0, 3, 4, 5, 8, 9, 10, 14, 15)$ using **3-input NOR gates**

$$f' = a'b'c'd + ab'cd + abc' + a'bc + a'cd'$$

⇒ two-level circuit:

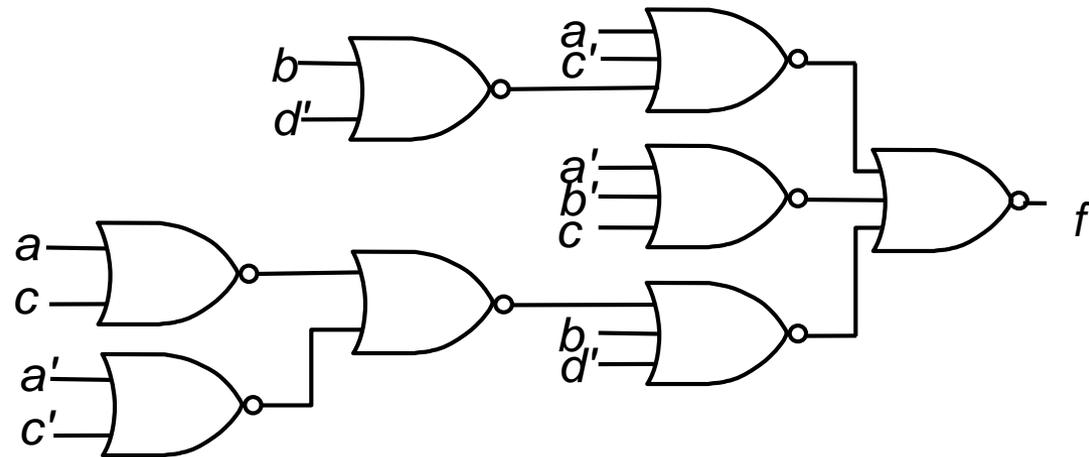
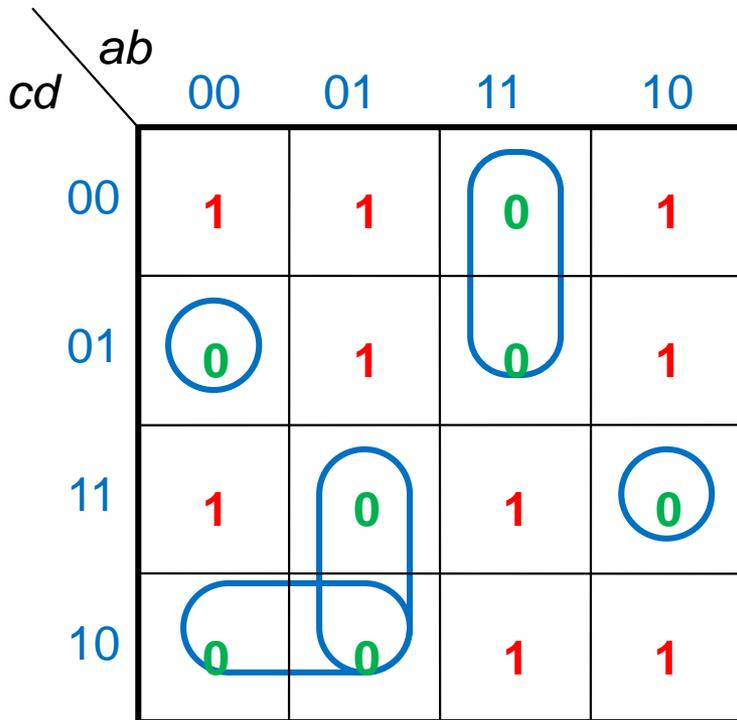
2 4-input AND (X) ⇒ factoring

3 3-input AND (O)

1 5-input OR (X) ⇒ sum of 3 items

$$\Rightarrow f' = b'd(a'c' + ac) + a'c(b + d') + abc'$$

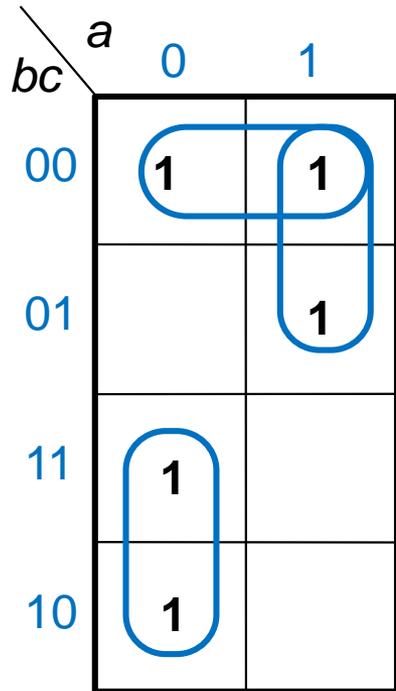
$$\Rightarrow f = [b + d' + (a + c)(a' + c)][a + c' + b'd][a' + b' + c]$$



f

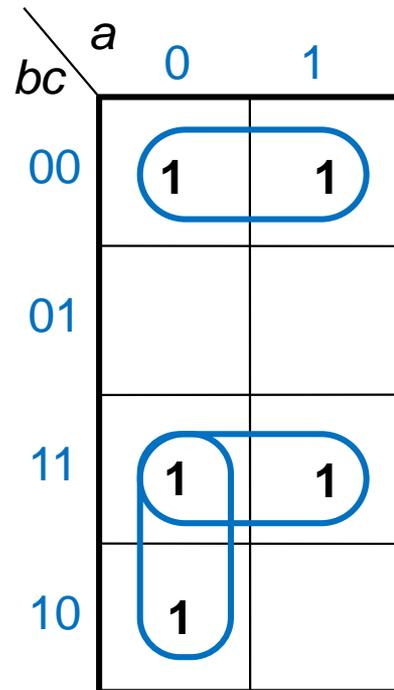
Example 2 (1/2)

- Realize the functions using only **2-input NAND** gates and **inverters**



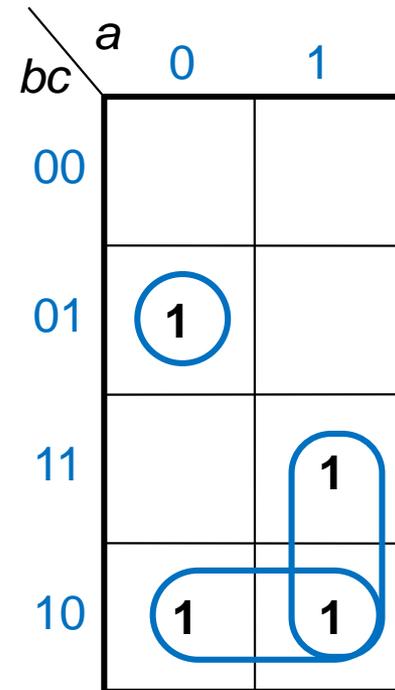
$$f_1 = \sum m(0,2,3,4,5)$$

$$f_1 = b'c' + ab' + a'b$$



$$f_2 = \sum m(0,2,3,4,7)$$

$$f_2 = b'c' + bc + a'b$$

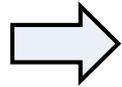


$$f_3 = \sum m(1,2,6,7)$$

$$f_3 = a'b'c + ab + bc'$$

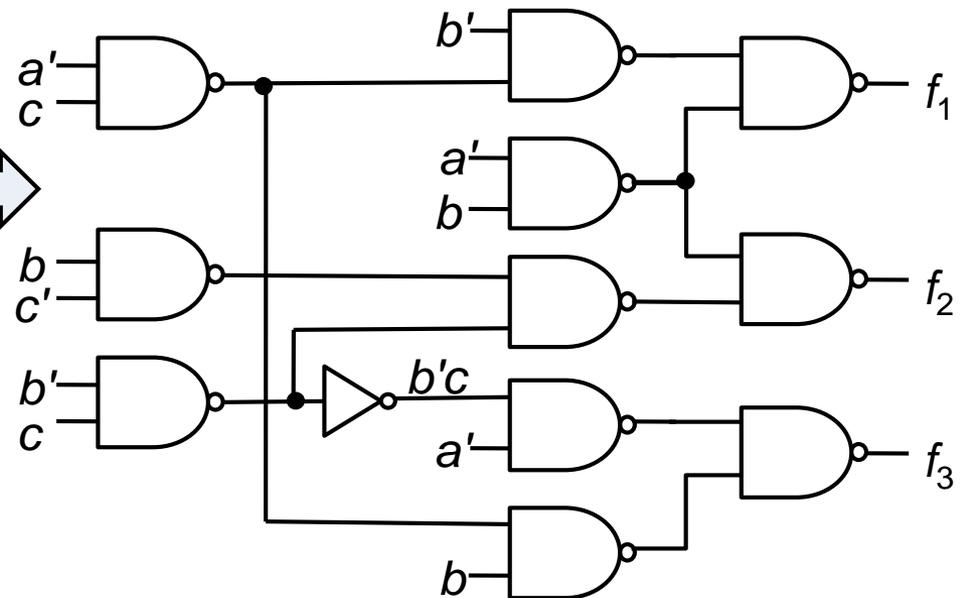
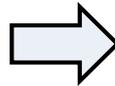
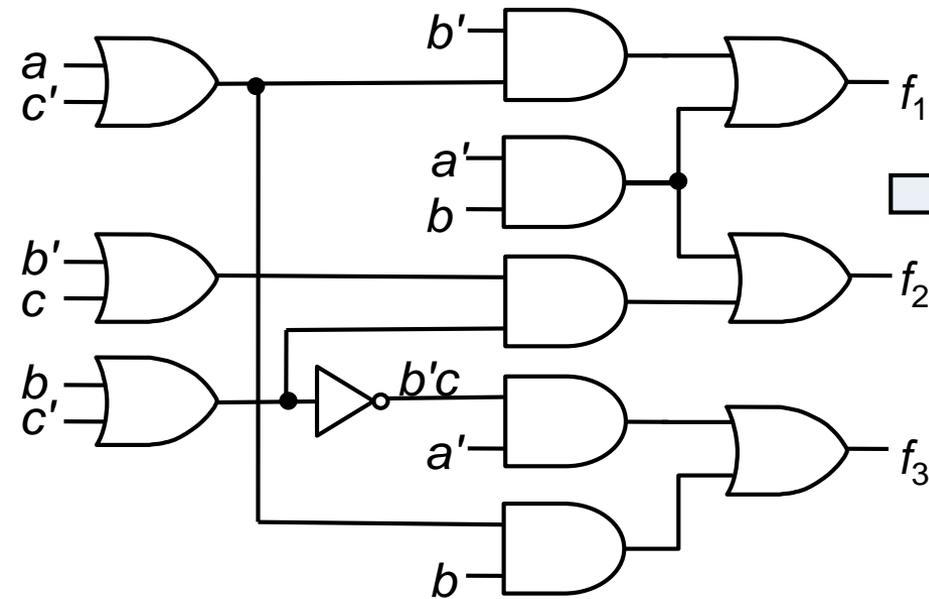
Example 2 (2/2)

$$\begin{aligned} f_1 &= b'c' + ab' + a'b \\ f_2 &= b'c' + bc + a'b \\ f_3 &= a'b'c + ab + bc' \end{aligned}$$



$$\begin{aligned} f_1 &= b'(a + c) + a'b \\ f_2 &= b'c' + b(a' + c) \quad \text{or} \quad f_2 = (b' + c)(b + c') + a'b \\ f_3 &= a'(b + c)' + b(a + c) \end{aligned}$$

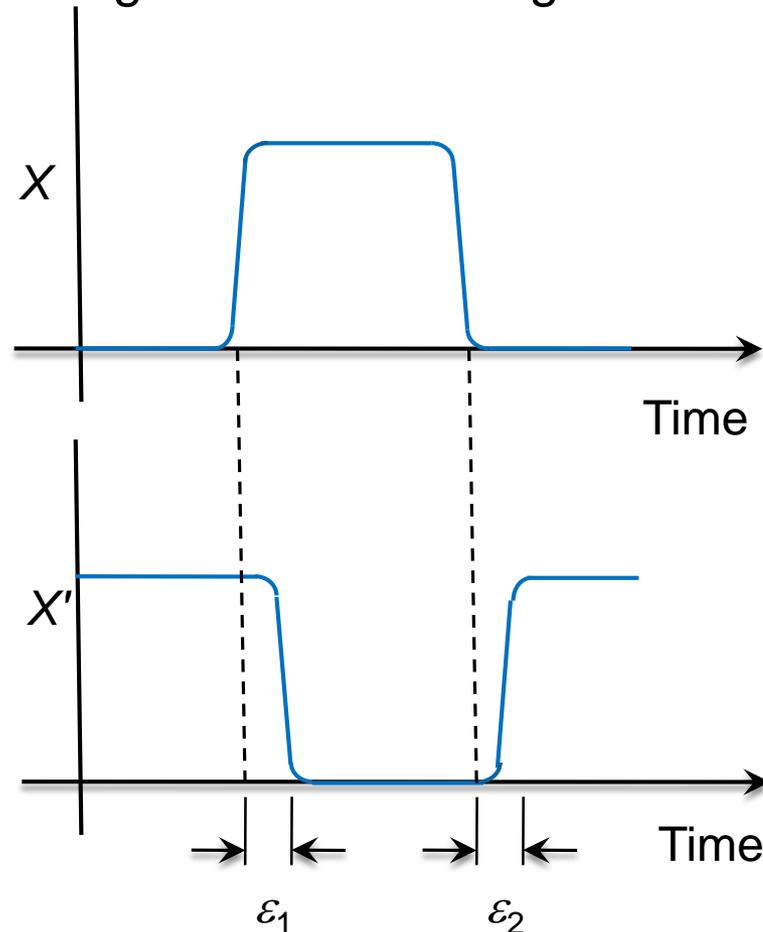
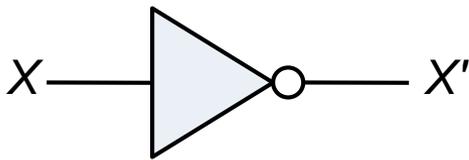
$a'b'c = a'(b'c) = a'(b + c)'$



8 Gate Delays and Timing Diagrams

Gate Delay

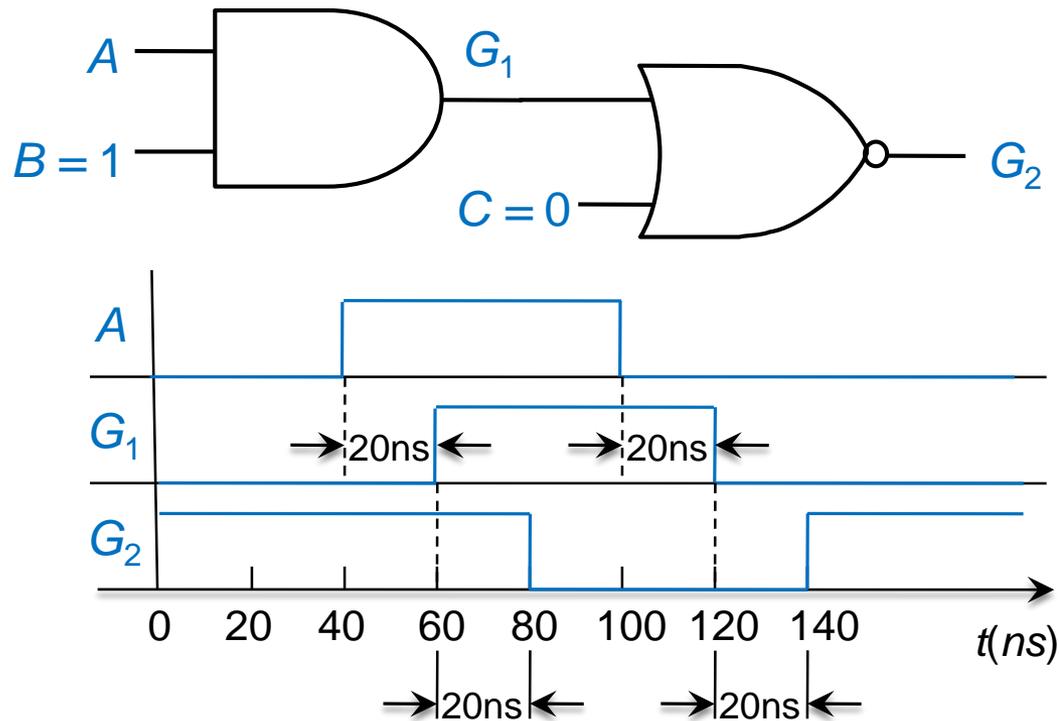
- **Propagation delay:** when the input to a logic gate is changed, the output takes some time to change value
 - ▣ Usually different for input rising $0 \rightarrow 1$ and falling $1 \rightarrow 0$



A few ns
(nanosecond = 10^{-9} sec)

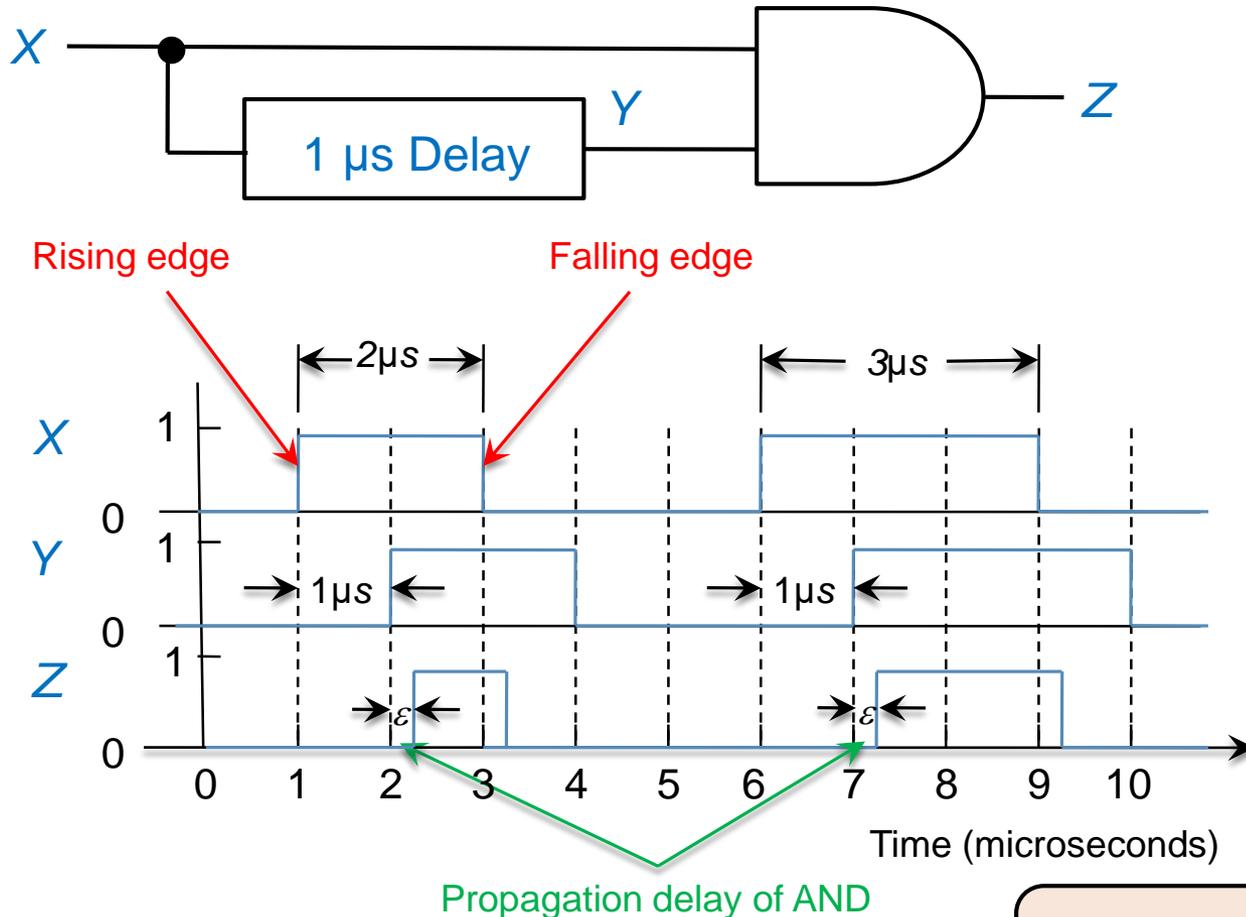
Sample Timing Diagram (1/2)

- e.g., assume each gate has a propagation delay of **20 ns**



Sample Timing Diagram (2/2)

- e.g., a circuit with a delay element



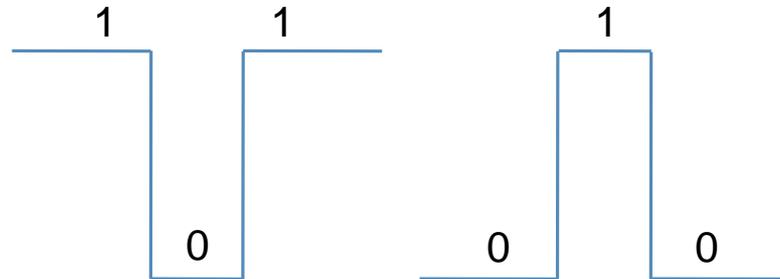
μs
(microsecond = 10^{-6} sec)

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Hazards

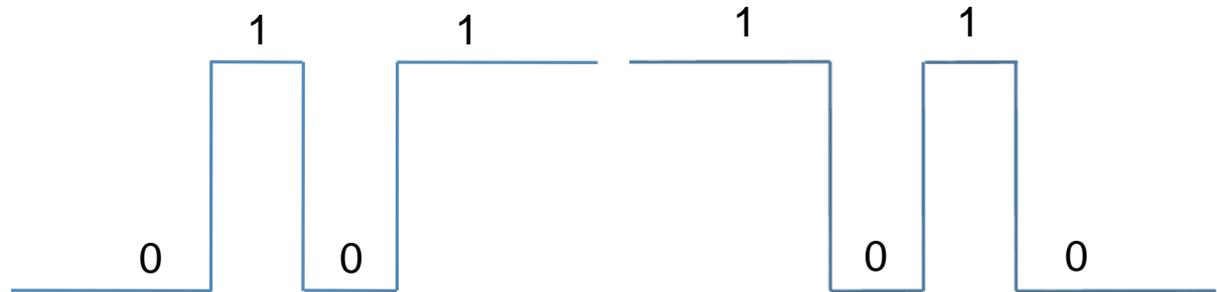
Hazards

- **Hazard**: unwanted switching transients appearing at the output when the input to a combinational ckt changes
 - Static 1-/0-hazard: output momentarily goes to 0/1 when it should remain a constant 1/0



(a) Static 1-hazard (b) Static 0-hazard

- Dynamic hazard: output change 3 or more times when the output changes from 0 to 1 (1 to 0)

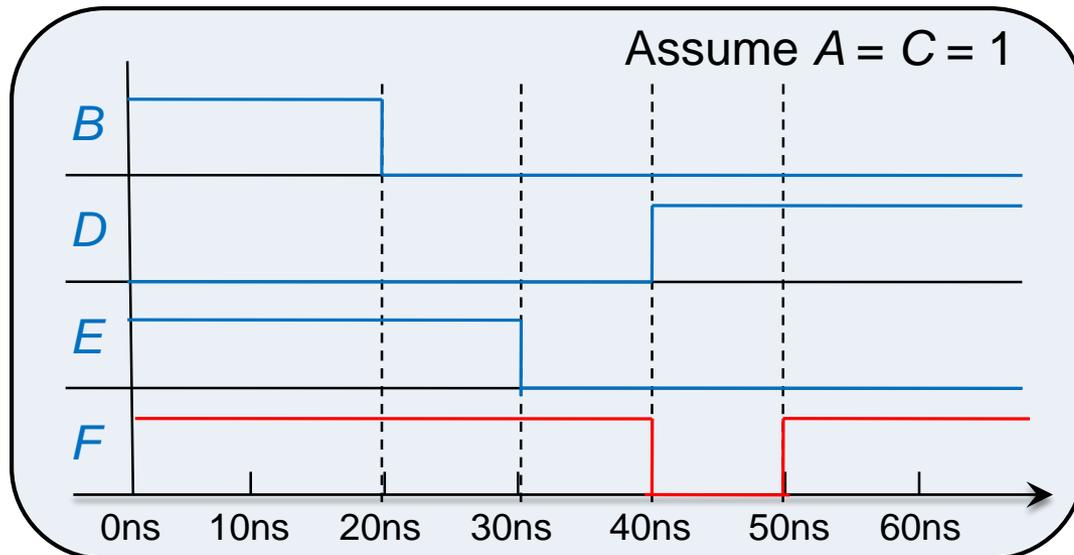
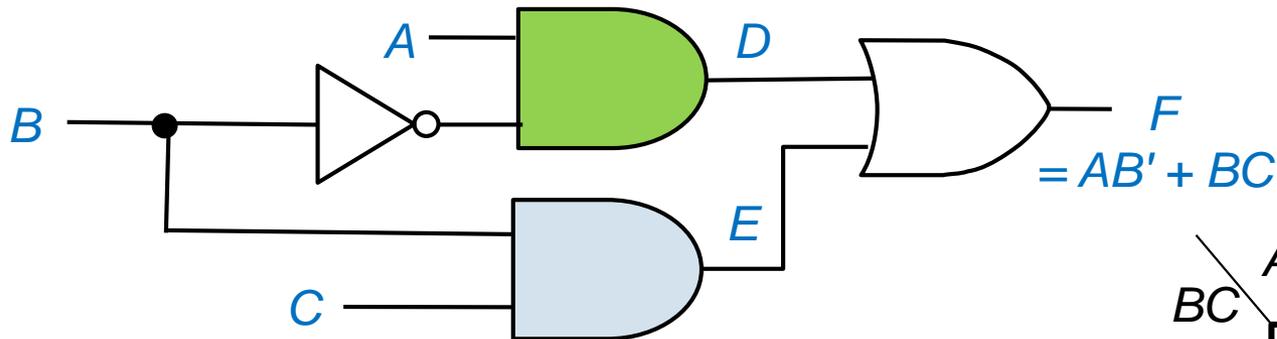


(c) Dynamic hazard

1-Hazards? When?

- e.g., assume each gate has a propagation delay of 10 ns

If $A = C = 1$, **static-1 hazard** occurs when B changes from 1 to 0



Combinational circuit design & simulation

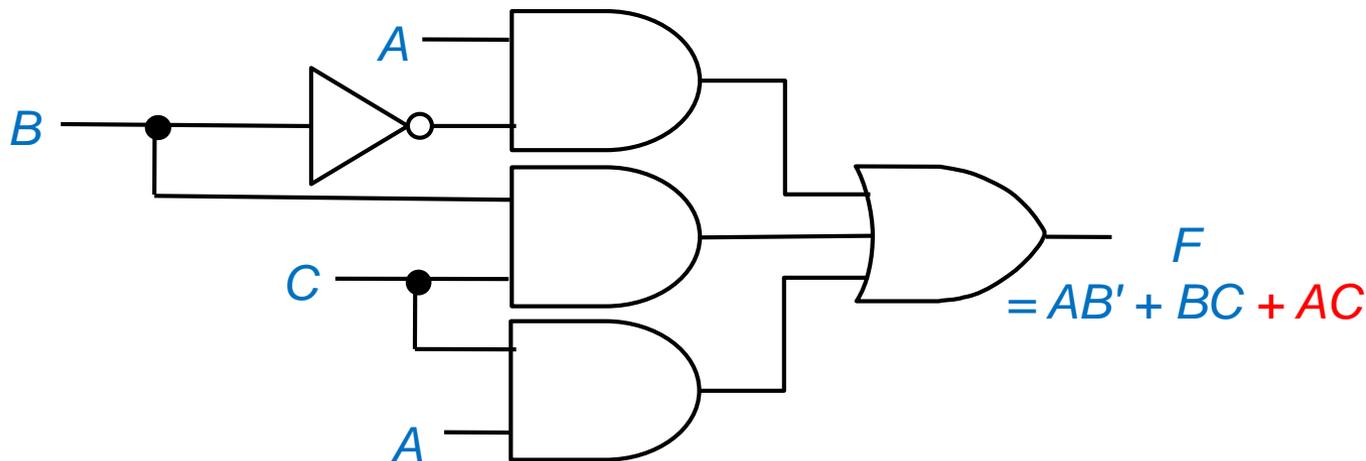
BC \ A	0	1
00	0	1
01	0	1
11	1	1
10	0	0

A green oval highlights the 1s in the 00 and 01 rows, labeled $AB'C$. A blue oval highlights the 1s in the 11 row, labeled ABC .

1-Hazard Removal for 2-Level AND/OR Circuit

- **Q: When?**
- **A: In Karnaugh maps, if any two adjacent 1's are not covered by the same loop, a 1-hazard exists for the transition between the two 1's**
- **Q: How to remove hazards?**
- **A: Add additional loops s.t. all adjacent 1's are covered by some loop**

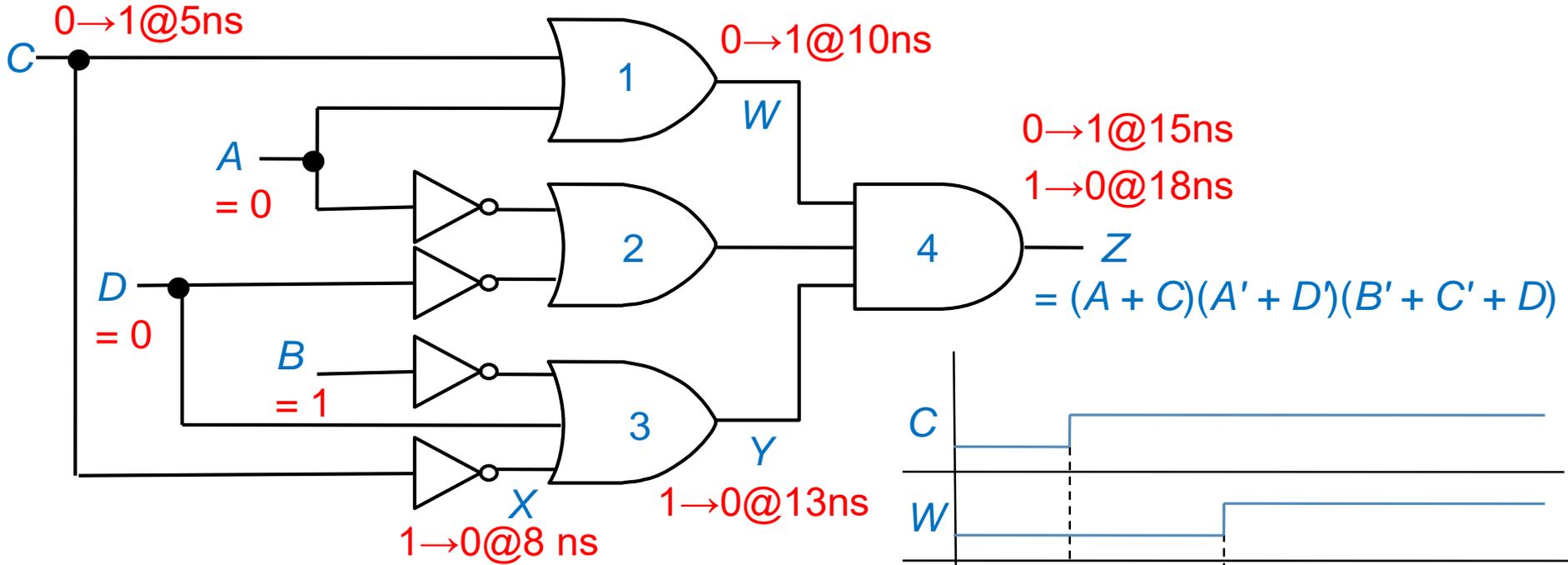
Consider only one-input change!



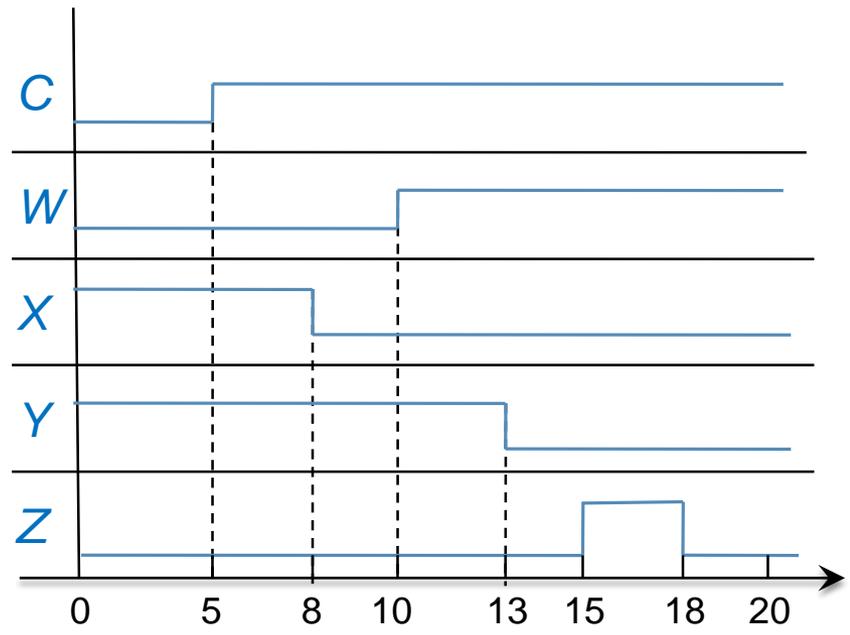
BC \ A	0	1
00	0	1
01	0	1
11	1	1
10	0	0

0-Hazards? When?

□ e.g., consider $A = 0, B = 1, D = 0, C$ changes from 0 to 1



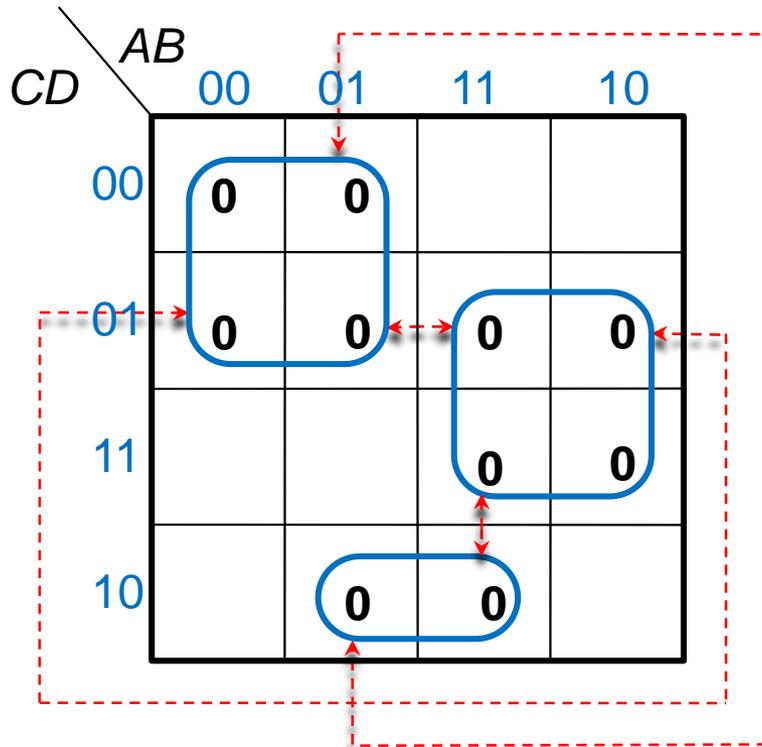
Propagation delay:
Inverter: 3 ns
AND/OR: 5 ns



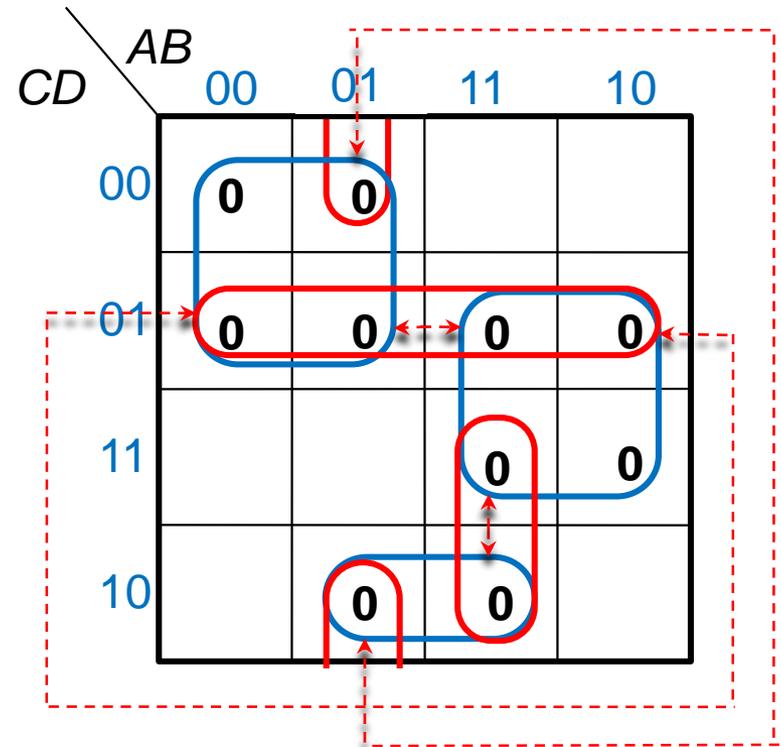
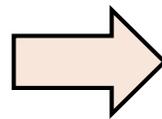
0-Hazard Removal for 2-Level AND/OR Circuit

- Similarly, add additional loops for adjacent 0's
- e.g.,

Consider only one-input change!



$$Z = (A + C)(A' + D')(B' + C' + D)$$



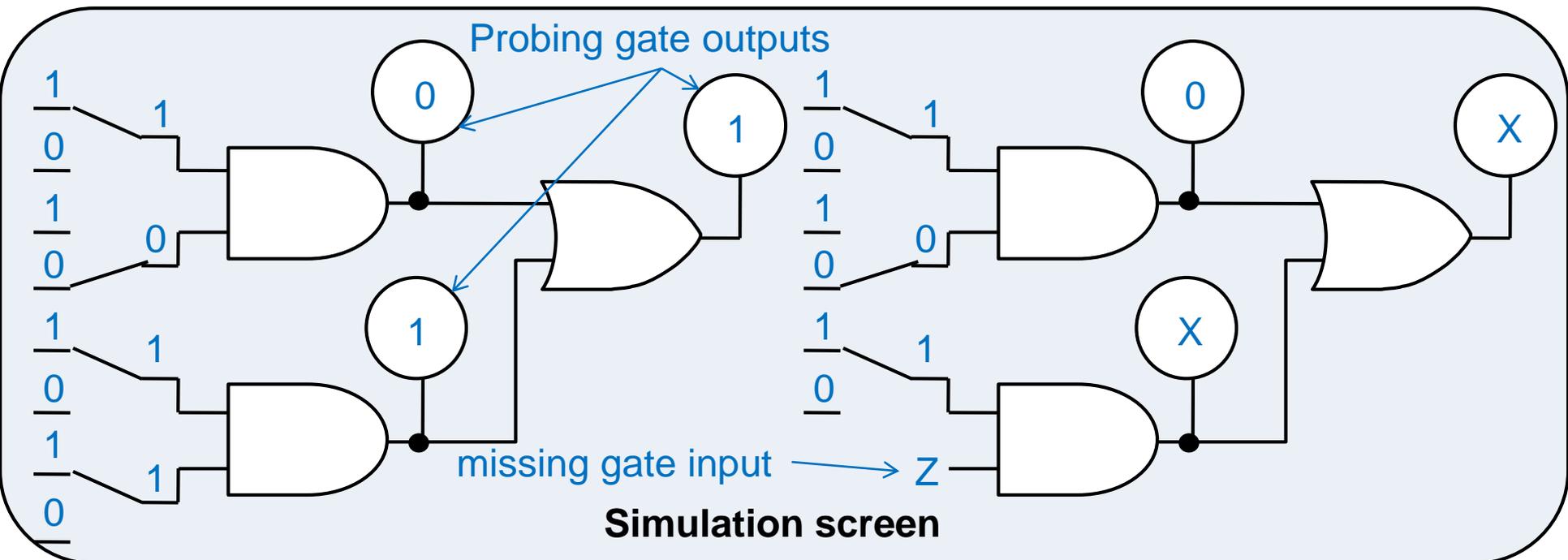
$$Z = (A + C)(A' + D')(B' + C' + D)(C + D')(A + B' + D)(A' + B' + C)$$

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Simulation

Simulation

- **Verify logic circuits by**
 - ▣ Actually building them
 - ▣ Simulating them on a computer
 - Computations are performed **stage-by-stage** from inputs
 - Evaluations are performed every time inputs change
- **4 logic values:** 0 (low), 1 (high), **X (unknown)**, **Z (high-Z/open ckt)**



4-Valued Logic Simulator

□ AND and OR functions for 4-valued simulation

•	0	1	X	Z	+	0	1	X	Z
0	0	0	0	0	0	0	1	X	X
1	0	1	X	X	1	1	1	1	1
X	0	X	X	X	X	X	1	X	X
Z	0	X	X	X	Z	X	1	X	X

□ Possible errors in designs

- In simulation
 - Incorrect design
 - Gates connected wrong
 - Wrong input signals to the circuits
- In real circuits
 - Defective gates
 - Defective connecting wires